Traffic Characteristics and Measurement: Internet Measurements

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Lecture Outline Outline

- 1. Why Capture Network Traffic?
- 2. Packet Capturing
- 3. Flow Capturing
- 4. Case-Study: The DDoSVax-Project
- 5. Related Student Thesis offers @ CSG

1. Why Capture Network Traffic?

- Accounting
- Network monitoring
- Forensics
- Research
- Intelligence
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Tasks Related to Data Capturing

- Obtaining/creating suitable sensors
- Sensor placement and operation
- Short-term storage at or close to sensor
- Transfer off-site
- (Long-term) storage
- Processing (libraries, infrastructure)
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2. Packet Capturing

Capturing alternatives

- Complete packets, i.e. complete Layer 3 payload
- Layer 3+4 packet headers (Typically first 48 bytes of Layer 3 packet)
- First 60 bytes of each Layer 3 packet

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Reasons to Capture Packets

- "Complete" network traffic, most accurate
- No aggregation or preprocessing needed
- Single packet timing and sizes
- Application layer data
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Reasons Not to Capture Packets

- Difficult/impossible for fast links
- Massive amount of data on fast links
 - difficult to store (short and long term)
 - difficult to transfer
 - difficult to process
- Data payloads often not needed
- May be illegal to do

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Packet Sensors

- Fast Ethernet (FE), 100Mb/s: Standard PCs with Linux, FreeBSD, ...
- Gigabit Ethernet (GbE): (Fast PC), network processor, special hardware Special software, tailored to the hardware
- Solution States (Constant States)
 Solution States (Constant States)
- Alternative for specific traffic: (More or less) transparent proxies

Sensor Placement I

- Needs to see all traffic to/from site of interest
- Needs to have data transfer resources
- May need to be invisible (intelligence operations)

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Sensor Placement Example



Valid Placements: A or (B and C and D)!

Sensor Placement II

- Capturing a link with bandwidth n gives two packet streams with bandwidth n!
- Capturing at a router with m links that have bandwidth n gives

2 * m * n

bandwidth to be captured!

Storage and Transfer

Example: Gigabit Ethernet (GbE), single bidirectional link

- Full traffic, worst case:
 200MB/s = 720GB/h = 17TB/day = 6.3PB/year
 Needs >= 2.5GbE dedicated link for reliable transfer
- Headers only: ?

Headers: Same worst-case (SYN-flooding, ICMP,...).

- PCI bus: 135 MB/s
- Harddisk/tape-drive: 50MB/s

About Time-Stamps

Timestamps are critical to correlate data from multiple sensors.

- May be needed to determine packet sequence
 Nanosecond accuracy may be needed!
- May have absolute time $\Rightarrow >= 32$ bit for seconds

May be up to 64 bit per timestamp.

Processing

- Real-time: Same problems as capturing and storage
- Connections need to be reconstructed
 May need partial network stack
- Payload processing has arbitrary complexity

Legal and Ethical Aspects

Disclaimer: I am no expert!

- Payloads may fall under privacy laws
 ⇒ Capturing and/or storage may be illegal
- Respect the privacy of individuals !/?
- Payloads may contain passwords, credit-card numbers, etc.
 - \Rightarrow Liability if misused (e.g. identity-theft) ?
- Criminal activity in payloads: Obligation to report?

Headers are far less problematic.

Summary for Packets

- Very accurate
- Difficult/expensive to capture, transfer, store, process
- Long-term continuous monitoring often infeasible
- Real-time monitoring difficult
- May cause legal problems

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3. Flow Traces

A *network flow* is an aggregated stream of packets from one source (IP, port) to one destination (IP, port).

- Addresses: IP-addresses, ports (TCP, UDP)
- Timestamp: First packet, last packet
- Counters: Bytes, packets
- Flags (TCP): SYN, FIN, RST, ...
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Limitations of Flow Capturing

Examples:

- Inaccurate/incomplete header information
- No payload information
- No packet sizes
- Maximum flow duration (e.g. 15 minutes)
- Maximum idle timeout (e.g. 30 seconds)
- Maximum data length (e.g. 4GiB)

Flow Data Format Alternatives

- NetFlow v5 (v7)
- NetFlow v9 (not yet implemented widely)
- IPFIX (still in definition)

Bi-directional unusual, Internet has asymmetric routing!

Example: NetFlow v5

- Available in current Cisco Routers
- Exports UDP packets from the routers
- 24 Byte packet header
- 48 Bytes per flow
- Grown historically

NetFlow v5 UDP Packet Header

struct netflow_v5_header {

- uint16_t version;
- uint16_t count;
- uint32_t SysUptime;
- uint32_t unix_secs;
- uint32_t unix_nsecs;
- uint32_t flow_sequence;
- uint8_t engine_type;
- uint8_t engine_id;

uint16_t reserved;

};

NetFlow v5 UDP packet Flow Record

struct netflow_v5_record {

uint32_t	addr;	uint32_t	dstaddr;
uint32_t	nexthop;	uint16_t	input;
uint16_t	output;	uint32_t	dPkts;
uint32_t	dOctets;	uint32_t	First;
uint32_t	Last;	uint16_t	port;
uint16_t	dstport;	uint8_t	pad1;
uint8_t	tcp_flags;	uint8_t	prot;
uint8_t	tos;	uint16_t	_as;
uint16_t	dst_as;	uint8_t	_mask;
uint8_t	dst_mask;	uint16_t	pad2;

Flow Sensors

- Typically router-integrated ("free")
- Export e.g. via UDP
- Export can be in dedicated link or within normal traffic
- Data-rate limited by sensor (limited buffer) \Rightarrow Data loss with too many short flows
- In fast networks sampling may be used

Transfer, Storage

- Typically feasible with commodity hardware
- Long-term storage needs tape/disk library
- Compression unproblematic

See case study for more information

4: Case-Study: DDoSVax Project

http://www.tik.ee.ethz.ch/~ddosvax/

- Collaboration between SWITCH (www.switch.ch) and ETH Zurich (www.ethz.ch)
- Aim (long-term): Analysis and countermeasures for DDoS-Attacks and Internet Worms
- Start: Begin of 2003
- Funded by SWITCH and the Swiss National Science Foundation

SWITCH

The Swiss Academic And Research Network

- .ch Registrar
- Links most (all?) Swiss Universities
- Connected to CERN
- Carried around 5% of all Swiss Internet traffic in 2003
- Around 60.000.000 flows/hour
- Around 300GB traffic/hour

SWITCH Peerings



SWITCH Traffic Map



NetFlow Data Usage at SWITCH

- Accounting
- Network load monitoring
- SWITCH-CERT, forensics
- DDoSVax (with ETH Zurich)

Transport: Over the normal network

NetFlow Data Flow



NetFlow Capturing

- One Perl-script per stream
- Data in one hour files
- Timestamps and src-IP in "stat" file

Critical: Linux socket buffers:

- Default: 64kB/128kB max.
- Maximal possible: 16MB
- We use 2MB (app-configured)
- 32 bit Linux: May scale up to 5MB/s per stream

Capturing Redundancy

- Worker / Supervisor (both demons)
- Super-Supervisor (cron job)
 For restart on reboot or supervisor crash
- Space for 10-15 hours of data

No hardware redundancy

Data Transfer to ETHZ

- Cron job, every 2 hours
- Single Perl script
- Transfer: scp (no compression, RC4)
- Remote deletion: ssh

No compression on ezmp2. (Some other Software running there) Bzip2 compression on ezmp2 would be possible!

Long-Term Storage Format

Full data since March 2003 Bzip2 compressed raw NetFlow V5 in one-hour files

- We need most data and precise timestamps
- We don't know what to throw away
- We have the space
- Preprocessing for specific work still possible

Latency: 5-10 minutes / hour of data

Computing Infrastructure

The "Scylla" Cluster
http://www.tik.ee.ethz.ch/~ddosvax/cluster/
Servers:

- aw3: Athlon XP 2200+, 600GB RAID5, GbE
- aw4: Dual Athlon MP 2800+, 800GB RAID5, GbE
- aw5: Athlon XP 2800+, 800GB RAID5, GbE

Nodes:

• 22 * Athlon XP 2800+, 120GB, GbE

Information somewhat outdated.

Infrastructure Cost (2004)

Hardware and full installation:

- aw3 (capturing): 1600 USD + 2 MD
- aw4 (dual CPU server): 2500 USD + 3 MD
- Cluster: 24.000 USD + 1MM
- Maintenance: 1-2 MD/month

Hidden cost: Computer room, network infrastructure, software development Scalability: Add 2*200GB HDD to each node \Rightarrow 8TB additional at 6000 USD

Lessons learned

Most important: KISS!

- Use scripting wherever possible
- Worker and Supervisor pairs are simpler
 ⇒ "crash" as error recovery model
- Cron as basic reliable execution service
- Email for notification: Do rate-limiting
- File-copy: Interlock and age check
- ssh, scp password-less (user key)
- Nothing needs to run as "root"!

Remarks on Software

- Linux is stable enough
- Linux is fast enough
- Linux Software RAID1/5 works well
- XFS has issues with Software RAID
- Perl is suitable for demons
- Python is suitable for demons

Remarks on Hardware

PC hardware works well, but:

- Get good quality components (PSUs!)
- Get good cooling (HDDs/CPUs)
- Do SMART monitoring
- Do regular complete surface scans
- Have cold spares handy
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Remarks on Linux Clusters

- Rackmount vs. "normal"
- Cooling / Power needs planning
- Gigabit Ethernet "star" topology is nice
- KVM not for all nodes needed
- FAI (Fully Automatic Installation) for installation
- Local Debian mirror
 - \Rightarrow 10 Min for complete reinstallation
- No global connectivity for the nodes
- Private addresses for the nodes

UPFrame

http://www.tik.ee.ethz.ch/~ddosvax/upframe/

- UDP plugin framework
- E.g. for online analysis of NetFlow data
- Can be used as traffic-shaper
- Robust: For experimental plugins

UPFrame Structure



Summary for Data and Infrastructure

- SWITCH is large enough and small enough
- No special hardware / software needed for capturing
- Long-term storage is unproblematic
- Linux can be used in the whole infrastructure
- Online processing is more difficult
- Simplicity and Reliability are the main issues
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The DDoSVax Dataset

- NetFlow v5 (converted from V7 by SWITCH)
- About 60.000.000 flows/hour
- Weekday: About 200k internal and 800k external IPs
- Unsampled
- Stored in full since March 2003

Flow Data Analysis by SWITCH

SWITCH-CERT: Short-term forensics (reduced)

- Single fast computer with hardware RAID-5
- No compression
- Sorted into minute (?) intervals
- Fast search with regular expressions
- Several weeks online
- No (?) long term storage

DDoSVax Offline Analysis

- E.g. for network/email worms
- Customised tools for some analyses
 - Single hour / prototyping: netflow_to_text and Perl
 - Days...weeks: From C-template
- Also other things: P2P, IRC, ...

Example: Blaster - Flows



Example: Blaster - Unique Sources



Example: Sobig



Sobig.F E-Mail Worm Propagation observed in the SWITCH Network

Example: MyDoom



Traffic Amount vs. Unique Sources

Traffic Amount:

- Easy to do
- Works reasonably well
- Sensitive to data generation problems
- Sensitive to observed network

Unique Sources:

- More complicated, more robust
- Weakly dependent on observed network
- Allows to get global picture

Analysis-tools: Scripting

"netflow_to_text"

- Takes one data file, outputs one line
- Well suited as "grep"/Perl input

Example:

TCP pr 111.131.210.8 si 1111.136.200.121 di 1264 sp 135 dp 48 le 1 pk 12:59:51.965 st 12:59:51.965 en 0.000 du

Analysis-tools: C

"Iterator template"

- Iterates over all records in a set of files
- Preprocesses timestamps, etc.
- Reading of input files encapsulated

Performance Issues

- S-10 minutes / hour of data bunzip2
- I/O limit at 10 cluster nodes reading from one NFS partition
- Memory limitations

5: Student Theses

http://www.tik.ee.ethz.ch/~ddosvax/sada/

Open topics for SA/DA/MA theses:

- Generic: Attack and Worm Outbreak Online Detection
- Visibility of World of Warcraft and Half-Life 2 traffic in NetFlow data
- Generic: Attack Detection and/or Signature Generation for Honeypots (Contact: B. Tellenbach, <betellen@tik.ee.ethz.ch>)

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Topics proposed by students are welcome!